

# Introduction to Computational Complexity

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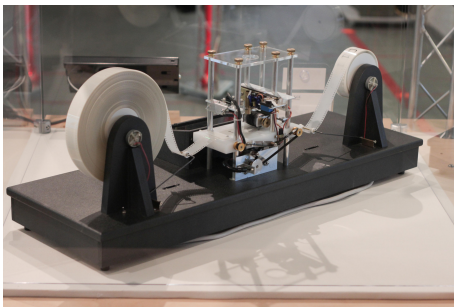
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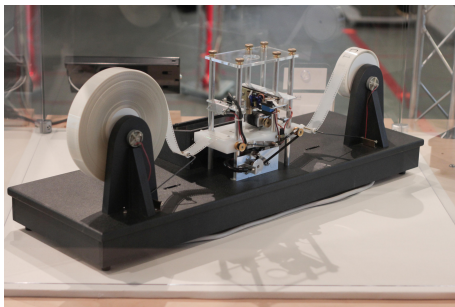
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# Computability

Central questions of theoretical computer science are connected to topics such as which problems are **computable** and **how efficiently**.



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- A fascinating thing about computational complexity is that we basically all agree on what can be computed.
- How efficiently? Much less clear.

# What does it mean to compute

By computing, we mean to evaluate a specific function  $f : \{0, 1\}^* \mapsto \{0, 1\}^*$  ( $S^* = \bigcup_{n \geq 0} S^n$  is the set of all finite strings over  $S$ ). This can be done by a procedure called **an algorithm**:

## Algorithm (informal)

An algorithm  $A$  computes a function  $f$  if:

- 1 It provides the output of  $f$  by following a **finite procedure** described by unambiguous elementary steps.
  - 2 Runs in a **finite number of steps** with no particular bound on the storage space used (it can always ask for more if needed).
- An algorithm is essentially what we understand as a **computable function**.
  - Notice that we do not speak about how efficient the computation of the function should be.

Power of a computational model:

- The sequence of steps that define the algorithm is executed by a **computational model** (e.g., mechanical machine, computer).
- But the computational model that manipulates with symbols should have a certain complexity (i.e., sufficiently powerful instruction set) to calculate (at least some) computable functions, right?
  - What operations are allowed? Random-access memory? Stack only? Conditional branching?
- What makes a computer a "universal" one?

# Turing machine is the universal model of computation

- **Turing machine (TM)**: the universal model of computation.
  - Abstract machine described by Alan Turing that reads symbols, changes its state, **rewrites symbols** on the tape, **moves the tape**.
  - Our computers  $\approx$  implementation of TM.

## Church-Turing conjecture (1936)

All computable functions are exactly Turing-computable (although not necessarily very efficiently).

- Not so obvious: we have examples that are known to be strictly less powerful: finite state machines, push-down automata, ...
- On the other hand, different computational models (e.g.,  $\lambda$ -calculus, TM with access to random bits) do not seem to be more powerful.
- We equate the intuitive concept of computable function with Turing-computable, which can be precisely defined.
- We can restrict our study of computational problems under the TM.

# Computational problems

- Computational problems can be seen as relations between the inputs (instances) and outputs (solutions).
- $x$  encoding of the instance,  $y$  encoding of the solution over some alphabet,  $S = \{0, 1\}$ ,  $S^* = \{\epsilon, 0, 1, 00, 01, 10, 11, 000, \dots\}$ .
- Let  $R(x, y) \subseteq S^* \times S^*$  be a relation. Each  $R$  defines a computational problem:

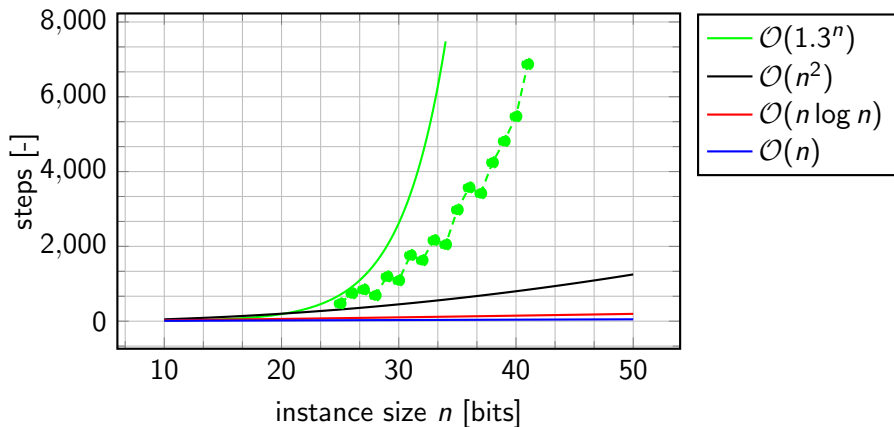
## Types of computational problems

- **Decision problems**
  - Given  $x$ , determine if there is  $y$  satisfying  $R(x, y)$ ?
- **Search problems**
  - Given  $x$ , find  $y$  such that  $R(x, y)$  or state it does not exist.
- **Optimization problems**
  - Given  $x$ , find  $y$  such that  $R(x, y)$  minimizing function  $c(x, y)$  or say no such  $y$  exists.
- **Function problems**
  - Compute value of  $f(x)$ .



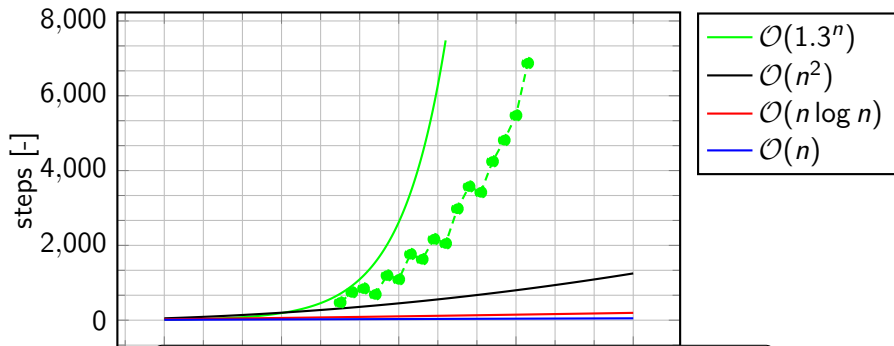
# Complexity of problems

- Not all  $R$  (computational problems) are equally difficult.
- We can measure the difficulty of the problem by the number of steps  $T(n)$  the best-known algorithm  $A$  on a TM needs to solve the problem  $R$  for a given input length  $n = |x|$  in **the worst-case**.



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**Computers are not to rescue:** electron-sized transistor, clock time  $\approx$  time of light between atoms, Earth-sized computer needs billions of years to brute force 300-bit solution.

# Problems, instances and algorithms: summary

- A computational **problem** is a relation over **instances** and **solutions**.
- To solve problems, we develop **algorithms** with certain **time complexity**.
  - Measured in terms of the worst-case number of steps  $T(n)$  over all instances of length  $n$ .
- The existence of an algorithm with given time complexity  $\mathcal{O}(T(n))$  is a witness of the **problem** being in certain **complexity class**.
  - People started to categorize problems into a taxonomy.
- It turned out that there is a fundamental barrier between the polynomially solvable problems and the others.
  - In practice, we usually get low-degree polynomial algorithms or exponential ones (or even worse).

It motivates us to study which problems fall into the "good" category and which fall into "naughty".

# Efficiently solvable: P

- We will use decision problems (yes/no) to demonstrate the most prominent complexity classes.
  - Similar can also be done for optimization problems, with a few definition adjustments.

## Definition: Class P

The set of problems that are **solvable** in a polynomial time.

- For every  $x \in \{0, 1\}^*$  they state if  $\exists y \in \{0, 1\}^* : R(x, y)$  or no.
- Admit  $\mathcal{O}(\text{poly}(n))$  algorithm, e.g.  $\mathcal{O}(n^2)$ ,  $\mathcal{O}(n \log n)$ .

## Examples

- Integer number problems: Addition, Multiplication, Primality test,...
- Graph problems: Topological Sorting, Minimum Spanning Tree, ...
- Miscellaneous problems: Discrete Fourier Transform, Linear Programming, ...

# Efficiently checkable: NP

Definition: Class NP (non-deterministic polynomial)

The set of decision problems whose solutions are **checkable** in a polynomial time.

- What do we mean by checkable?

# Efficiently checkable: NP

## Definition: Class NP (non-deterministic polynomial)

The set of decision problems whose solutions are **checkable** in a polynomial time.

- What do we mean by checkable?
- **YES**-instances have so-called **polynomial certificates** (or witnesses, proofs): e.g.,  $|y| \leq \text{poly}(|x|)$ .
- Given certificate  $y$ , one can in polynomial time verify that indeed  $(x, y) \in R$ .
- We do not know whether they admit a  $\mathcal{O}(\text{poly}(n))$  algorithm, but we know that their solution is verifiable by  $\mathcal{O}(\text{poly}(n))$  algorithm.  $\approx$  there is an algorithm that solves the problem in polytime given the certificate  $y$ .

## Examples

- Integer number problems: Sudoku, Knapsack, 2-Partition, ...
- Graph problems: Travelling Salesman,  $k$ -coloring, ...
- Miscellaneous problems: Linear equations with absolute values, Control theory: constrained state-space feedback

# SubsetSum is in NP

## Definition: SubsetSum problem

- **Instance:** A (multi)-set of  $n$  non-negative integers  $A = \{a_1, \dots, a_n\}$  and a non-negative integer  $W$ .
- **Decision:** Is there a subset  $S \subseteq A$  such that  $\sum_{a_i \in S} a_i = W$ ?

## Example 1: YES-instance

$A = \{1, 1, 2, 3, 7, 9\}$ ,  $W = 6$ . The answer is **YES**:  $S = \{1, 2, 3\}$ .

- I claim  $(A, W)$  is YES-instance. This is a poly-sized proof:  $S$ .

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## Example 2: NO instance

$A = \{3, 5, 5, 6, 8, 10\}$ ,  $W = 25$ . The answer is **NO**.

- I claim  $(A, W)$  is NO-instance. I do not have a short proof.

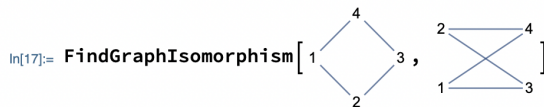
Generally, short certificates (proofs) of **NO**-instances may not exist.



# Graph Isomorphism is in NP

## Example: Graph Isomorphism problem

- Given two graphs  $G$  and  $H$ , decide if  $G$  is the same as  $H$  up to the vertex labelling.
- A  $\mathcal{O}(\text{poly}(n))$  algorithm is not known.
- We have an algorithm by Babai (2015) that runs in  $\mathcal{O}(\exp(\log(n)^c))$  for some constant  $c > 1$ , i.e., a quasipolynomial, grows smaller than an exponential.
- But its solution is checkable in a poly-time.



Out[17]= { <| 1 → 1, 2 → 3, 3 → 2, 4 → 4 |> }

- Remark:** in contrast to NP, problems in P have polynomial certificates even for **NO**-instances.

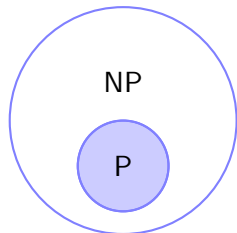
# What is non-deterministic about NP?

Why NP means "non-deterministic polynomial"?

- Connected with an alternative computational model, the so-called non-deterministic Turing Machine (TM).
- This abstract computational model explores all branches in your algorithm in parallel.
- This is an alternative description of class NP: the set of decision problems for which there is an algorithm that solves it in a polynomial time on a non-deterministic TM computational model.
- Useful for theoretical analysis, nobody knows how to build it physically (in contrast to the deterministic TM).
  - Quantum computers are not believed to be equivalent to non-deterministic TMs.

## How important is P vs NP question?

- At least \$1.000.000 important.
- Clay Math Institute's Millennium problems:
  - Solution smoothness of Navier–Stokes Equation
  - Poincaré Conjecture (solved)
  - Riemann Hypothesis
  - **P vs NP problem**
  - ...
- P vs NP question has wide implications to the world outside of CS: class P exactly corresponds to dynamical systems described by ODEs with polynomial RHS under a poly-length simulation (connection to control theory).



# P vs NP question

**Common belief is:**

Conjecture

$$P \neq NP.$$

- Likely we are not in a position to resolve this question within the next 20 years.

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<sup>1</sup><https://www.cs.umd.edu/users/gasarch/BLOGPAPERS/pollpaper3.pdf>

<sup>2</sup>[https://youtu.be/pQsdygaYcE4?si=N\\_22d0eZyHLeUngt](https://youtu.be/pQsdygaYcE4?si=N_22d0eZyHLeUngt)

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- But, we can run a survey:

Year	2002	2012	2019
Thinks $P \neq NP$	61%	82%	88%

Table: William Gasarch's survey on P vs NP<sup>1</sup>.

- See nice explanatory video on P vs NP<sup>2</sup>.

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# P vs NP question

Common belief is:

Conjecture

- Likely we are not in 20 years.
- But, we can run a simulation



THE CLASSIC WORK  
NEWLY UPDATED AND REVISED

## The Art of Computer Programming

VOLUME I  
Fundamental Algorithms  
Third Edition

DONALD E. KNUTH

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Thinks $P \neq NP$	61%	82%	88%

Table: William Gasarch's survey on P vs NP<sup>1</sup>.

- See nice explanatory video on P vs NP<sup>2</sup>.
- Donald Knuth is one of the great proponents of  $P = NP$ .

<sup>1</sup><https://www.cs.umd.edu/users/gasarch/BLOGPAPERS/pollpaper3.pdf>

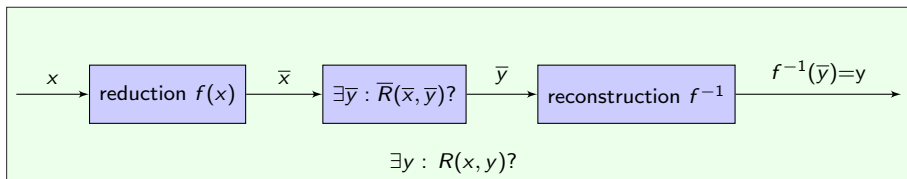
<sup>2</sup>[https://youtu.be/pQsdygaYcE4?si=N\\_22d0eZyHLeUngt](https://youtu.be/pQsdygaYcE4?si=N_22d0eZyHLeUngt)

# Polynomial reductions

**Problem reductions** are one of the greatest inventions in computer science.

**Motto:** *My problems are your problems.*

- Solving a new problem  $R(x, y)$  via existing problem  $\bar{R}(\bar{x}, \bar{y})$ :

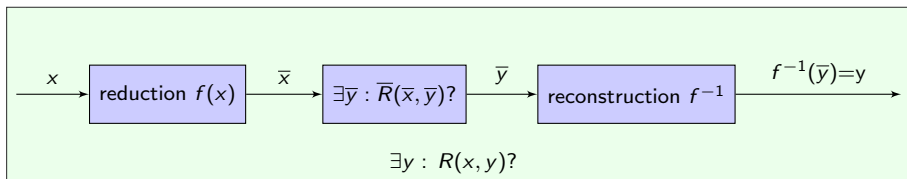


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Namely, we will be interested in **polynomial-time reductions**.

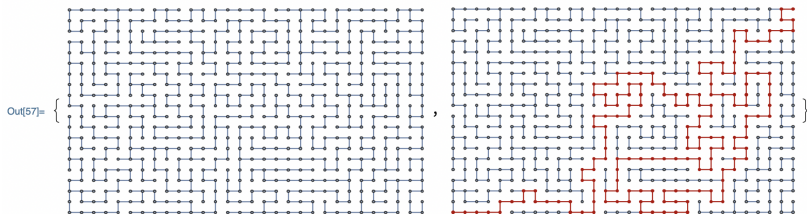
- $f$  and  $f^{-1}$  runs in a polynomial time
- Preserve membership in classes P and NP:  $poly(poly(n)) \in \mathcal{O}(poly(n))$ .
- Useful from the practical standpoint.



# You have used polynomial reductions before

- Path with the minimum number of edges, but you only have Dijkstra.

```
{m, HighlightGraph[m, PathGraph@FindShortestPath[m, First@VertexList[m], Last@VertexList[m]]]}
```



- This is a polynomial reduction.

# Some the reductions connect different worlds

- More examples: 3CNF-SAT  $\leq_P$  SubsetSum ( $R(x, y) \leq_P \bar{R}(\bar{x}, \bar{y})$ )

## Definition: 3CNF-SAT Problem

- **Instance:** A propositional formula in conjunction normal form with clauses with 3 literals, e.g.,  $\phi = (x \vee \neg y \vee z) \wedge \dots \wedge (\neg x \vee u \vee \neg v)$ .
- **Decision:** Is the formula  $\phi$  satisfiable?

## Example

$$\phi = (x \vee y \vee z) \wedge (x \vee y \vee \neg z) \wedge (x \vee \neg y \vee z) \wedge (x \vee \neg y \vee \neg z) \wedge \\ (\neg x \vee y \vee z) \wedge (\neg x \vee y \vee \neg z) \wedge (\neg x \vee \neg y \vee z) \wedge (\neg x \vee \neg y \vee \neg z)$$

The answer is **NO**.

# 3CNF-SAT to SubsetSum

- More examples:  $3\text{CNF-SAT} \leq_P \text{SubsetSum}$  ( $R(x, y) \leq_P \bar{R}(\bar{x}, \bar{y})$ )

## Definition: SubsetSum problem

- **Instance:** A (multi)-set of  $n$  non-negative integers  $A = \{a_1, \dots, a_n\}$  and a non-negative integer  $W$ .
- **Decision:** Is there a subset  $S \subseteq A$  such that  $\sum_{a_i \in S} a_i = W$ ?

## Example

$A = \{1, 1, 2, 3, 7, 9\}$ ,  $W = 6$ . Answer is **YES**:  $S = \{1, 2, 3\}$ .

- How can we use a number counting problem to solve a logic problem?  
These are different beasts.

# Example: 3CNF-SAT to SubsetSum

$$\phi = \underbrace{(\neg x \vee y \vee z)}_{C_1} \wedge \underbrace{(x \vee \neg y \vee z)}_{C_2} \wedge \underbrace{(\neg x \vee \neg y \vee \neg z)}_{C_3}$$

	x	y	z	C <sub>1</sub>	C <sub>2</sub>	C <sub>3</sub>	
x	1	0	0	0	1	0	a <sub>1</sub>
¬x	1	0	0	1	0	1	a <sub>2</sub>
y	0	1	0	1	0	0	a <sub>3</sub>
¬y	0	1	0	0	1	1	a <sub>4</sub>
z	0	0	1	1	1	0	a <sub>5</sub>
¬z	0	0	1	0	0	1	a <sub>6</sub>
	0	0	0	1	0	0	a <sub>7</sub>
	0	0	0	2	0	0	a <sub>8</sub>
	0	0	0	0	1	0	a <sub>9</sub>
	0	0	0	0	2	0	a <sub>10</sub>
	0	0	0	0	0	1	a <sub>11</sub>
	0	0	0	0	0	2	a <sub>12</sub>
	1	1	1	4	4	4	W

- Notice that no carry-overs are happening.
- **Homework:** does this work for  $k$ CNF-SAT ( $k$  literals in each clause)?

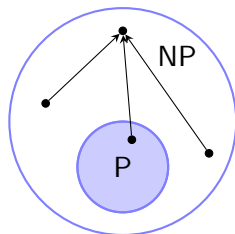
# Complete problems: NP-complete

The idea of reductions can be used to identify so-called **complete problems** for the class.

## Definition: NP-complete class

Problem  $R$  is NP-complete if  $R \in \text{NP}$  (i.e., efficiently checkable) and for every problem  $A$  :

$$\forall A \in \text{NP} : A \triangleleft_P R \text{ (i.e., acts as a solver).}$$



- The meaning of an NP-complete problem is that it represents a "universal" problem for NP class (can be used to solve all problems in NP).
- The first NP-complete problem was discovered by Cook (1971):
  - Proof: non-deterministic TM  $\triangleleft_P$  CNF-SAT.
  - Hence, CNF-SAT acts as a solver for class NP.
- Nowadays, we know thousands of NP-complete problems.

# Example: CNF-SAT

Problem reductions are not particularly useful if they do not run in a polynomial time.

## $k$ CNF-SAT Problem

- **Instance:** A propositional formula in conjunction normal form, e.g.,  
 $\phi = (x \vee \neg y \vee z) \wedge \dots \wedge (\neg x \vee u \vee \neg v)$ .
- **Decision:** Is the formula  $\phi$  satisfiable?

## $k$ DNF-SAT Problem

- **Instance:** A propositional formula in **disjunctive normal form**, e.g.,  
 $\phi = (x \wedge \neg y \wedge z) \vee \dots \vee (\neg x \wedge u \wedge \neg v)$ .
- **Decision:** Is the formula  $\phi$  satisfiable?

## Theorem

$k$ CNF-SAT is in NP-complete (NTM reduces to poly-sized CNF formula).  
 $k$ DNF-SAT is in P (easy algorithm).

**Reduction idea:** We have learned in TGR and LPS courses how to convert CNFs to DNFs (disjunctive normal form), and we know that DNF-SAT is solvable in a polynomial time (how?). So let's try

$$k\text{CNF-SAT} \leq_P k\text{DNF-SAT}.$$

# Example: CNF-SAT

```
In[12]:= cnf = (x1 ∨ y1) ∧ (x2 ∨ y2) ∧ (x3 ∨ y3) ∧ (x4 ∨ y4) ∧ (x5 ∨ y5) ∧ (x6 ∨ y6) ∧ (x7 ∨ y7);  
BooleanConvert[cnf] // TraditionalForm
```







# Example: CNF-SAT

Perhaps, if the DNF reduction would be done in a more sophisticated way:

```
in[16]:= BooleanMinimize[BooleanConvert[cnf]] // TraditionalForm
```

```
]]
```



## Takeaways:

- The above example shows an exponential explosion of the resulting DNF formula.
- Unfortunately, we do not know how to convert, in general, every CNF formula to DNF in a polynomial time.
- But reductions in the opposite direction, i.e., something  $\leq_P$  CNF-SAT, are in fact very useful:
  - formal verification: some states are not reachable within any  $k$  steps
  - proof checking: Keller's conjecture<sup>3</sup>
  - graph coloring, ...
- CNF-SAT is both theoretically (a universal NP problem) and practically (existence of solvers) appealing.

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<sup>3</sup><https://www.quantamagazine.org/computer-search-settles-90-year-old-math-problem-20200819/>

# NP-complete: summary

## NP-complete class summary:

- The set of universal (most difficult) problems for class NP.
- All algorithms we know for NP-complete problems have complexity above  $\mathcal{O}(\text{poly}(n))$ .
  - e.g., CNF-SAT algorithm is  $\mathcal{O}(1.308^n) \approx \mathcal{O}(2^{0.387n})$
- Solving efficiently one of the thousands known NP-complete problems would mean  $P = NP$ .
  - Hence, if your problem is NP-complete do not hope for a poly-time algorithm.

## NP problems suspected **not being** NP-complete and not in P

- Graph Isomorphism (GI)
  - We know a subexponential algorithm, but still above a polynomial complexity.
- Integer Factoring
- Computing VC (Vapnik-Chervonenkis) dimension

# NP-complete: summary

## NP-complete class summary:

- The set of universal (most)
- All algorithms we know for these problems are above  $\mathcal{O}(\text{poly}(n))$ .
  - e.g., CNF-SAT algorithm
- Solving efficiently one of these problems would mean  $P = NP$ .
  - Hence, if your problem is



**Complexity sandwich:** But can it be filled with natural ingredients?

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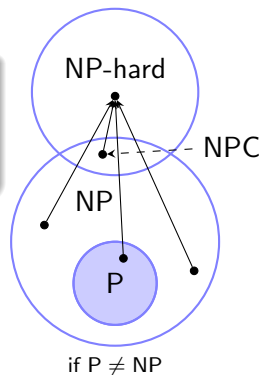
# Beyond NP-complete: NP-hard class

## Definition: NP-hard class

Problem  $R$  is NP-hard if for every problem  $A$

$$\forall A \in \text{NP} : A \leq_P R.$$

- sets a **lower bound** on the complexity of the problem (acts as a solver for class NP)
- the difference from NP-complete is that  $R$  can be much harder (does not have to be in NP)



## Examples

- every NP-complete decision problem
- optimization variants of NP-complete decision problems
- Quantified Boolean formula satisfiability:  $\forall x_1 \exists x_2 \forall x_3, \dots : f(x_1, x_2, x_3, \dots)$

## The main takeaways:

- Turing machine is the universal model of computation
  - Gives us a formal way studying and categorizing problems according to their complexity.
- Easy problems (P) vs. hard problems (NP-complete, NP-hard):
  - Easily solvable vs. easily checkable vs. just hard problems
  - More complexity classes live in Complexity ZOO:  
<https://complexityzoo.net/>.
- Polynomial reductions:
  - Using somebody else's problem to solve your problems.

- A. Borodin, Toronto Uni, CSC373<sup>4</sup>
- W. Gasarch: P vs NP survey<sup>5</sup>
- S. Aaronson: P vs NP status <sup>6</sup>
- L. Babai: Quasipoly algorithm for GI<sup>7</sup>
- Bournez et al.: Polynomial Time Corresponds to Solutions of Polynomial Ordinary Differential Equations of Polynomial Length<sup>8</sup>

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<sup>4</sup><https://www.cs.toronto.edu/~bor/373s13/>

<sup>5</sup><https://www.cs.umd.edu/users/gasarch/BLOGPAPERS/pollpaper3.pdf>

<sup>6</sup><https://www.scottaaronson.com/papers/pnp.pdf>

<sup>7</sup><https://arxiv.org/abs/1512.03547>

<sup>8</sup><https://arxiv.org/abs/1601.05360>